Lexicographically Fair Learning: Algorithms and Generalization

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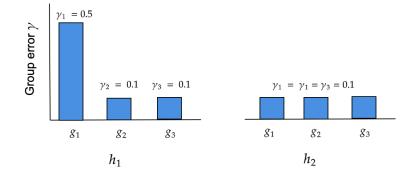
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We want our algorithms to treat different groups of people equitably.

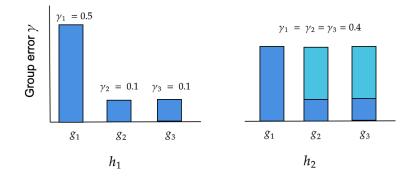
A Group Fairness Definition: Equality of Group Errors

"The algorithm should make the same number of mistakes on all groups."



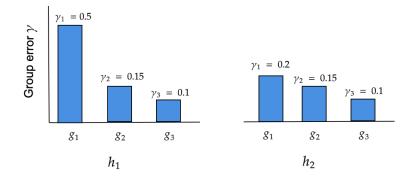
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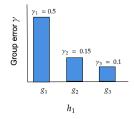
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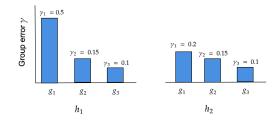


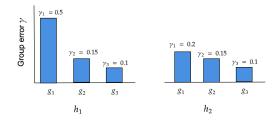
Alternative Group Fairness Definition: Minimax Group Fairness

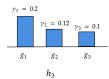
"The number of errors made on the worst-off group should be minimized."

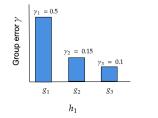


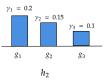


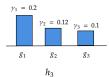


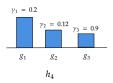




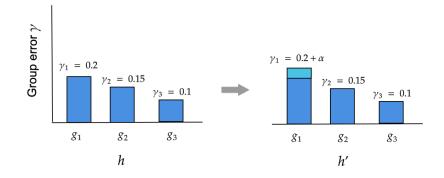




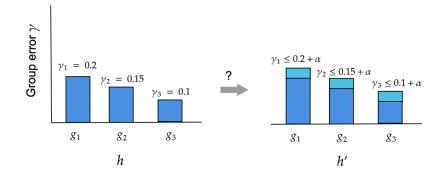


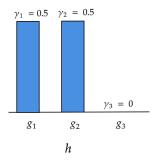


We can only efficiently get approximate minmax-fair solutions.

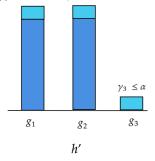


How do we generalize this to the lexifair setting?



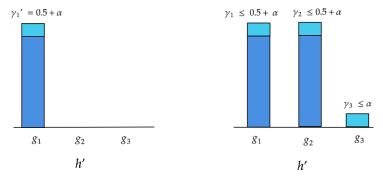


 $\gamma_1 \leq 0.5 + \alpha \qquad \gamma_2 \leq 0.5 + \alpha$



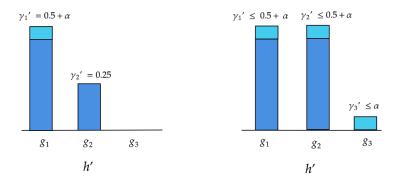
True lexifair solution

Approximate lexifair solution



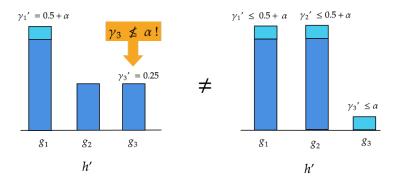
Approximate minimax solution

Approximate lexifair solution



Approximate min of top 2 group errors

Approximate lexifair solution



Approximate min of top 3 group errors

Approximate lexifair solution

Definition (Approximate Lexicographic Fairness)

Let $1 \leq \ell \leq K$ and $\alpha \geq 0$. Let $\vec{\epsilon} = (\epsilon_1, \epsilon_2, \dots, \epsilon_\ell)$, and define

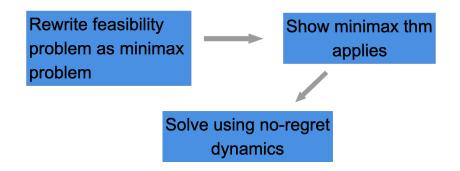
 $\mathcal{H}_{(0)}^{\vec{\epsilon}} \triangleq \text{ the entire model class } \mathcal{H},$ $\mathcal{H}_{(j)}^{\vec{\epsilon}} \triangleq \text{ models in } H_{j-1}^{\vec{\epsilon}} \text{ that have the smallest}$ *j*th group error rate up to an ϵ_i approximation.

A model *h* satisfies (ℓ, α) -lexicographic fairness ("lexifairness") if $h \in \mathcal{H}_{(\ell)}^{\vec{\epsilon}}$ for some $\vec{\epsilon}$ that is component-wise less than α .

- A constraint on the *highest* error amongst all groups, which arises in defining minimax error, is convex, and hence amenable to algorithmic optimization.
- However, naive specifications of lexifairness involve constraints on the second highest group errors, the third highest group errors, and more generally *k*th highest errors.
- These are non-convex constraints when taken in isolation.
- We get around this by replacing constraints on the *k*'th highest error groups with constraints on the *sums* of all *k*-tuples of group errors.

- Define a stable and convex version of approximate lexifairness.
- Derive oracle-efficient algorithms for finding approximately lexifair solutions.
- Show that when the underlying empirical risk minimization problem absent fairness constraints is convex, our algorithms are provably efficient.
- Show that approximate lexifairness generalizes: approximate lexifairness on the training sample implies approximate lexifairness on the true distribution w.h.p.

Oracle-efficient algorithms to achieve approximate lexifairness



- In regression setting, learner plays Online Projected Gradient Descent.
- In classification setting, learner plays Follow-the-Perturbed-Leader.

Algorithmic Formulation

- Our approach to find lexifair models is to **recursively** find the minimax (over sums of group error rates) rates
- Our algorithms return a model achieving those minimax rates, and hence that model will be lexifair.
- At level j, in an inductive fashion, we are given the minimax rates $\eta_1, \ldots, \eta_{j-1}$ from previous rounds, and we want to estimate η_j
- Can then dictate that every sum of j group error rates is at most η_i
- Writing the Lagrangian of this linear program

Let $L_{i_r}(h)$ indicate the loss incurred by the model h on the i_r 'th group. Then the Lagrangian for this linear program can be written as

$$\mathcal{L}_{j}((h,\eta_{j}),\lambda) = \eta_{j} + \sum_{r=1}^{j} \sum_{\{i_{1},\dots,i_{r}\}\subseteq[K]} \lambda_{\{i_{1},i_{2},\dots,i_{r}\}} \cdot (L_{i_{1}}(h) + \dots + L_{i_{r}}(h) - \eta_{r})$$
(1)

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Algorithmic Formulation: Two Player Zero-Sum Game

Can find a minimax solution for this Lagrangian with a zero-sum game between a (L)earner and a (A)uditor:



- At each round t, there is a weighting over groups determined by A
- L (best) responds by computing model *h*_t to minimize the weighted prediction error
- A updates group weights using online projected gradient descent with respect to group errors achieved by h_t
- L's final model M is uniform distribution over all of h_t 's produced

ALGORITHM 1: LexiFairReg: Finding a Lexifair Regression Model

Input: $S = \bigcup_{k=1}^{K} G_k$ data set consisting of K groups, (ℓ, α) desired fairness parameters, loss function parameters L_M

for
$$j = 1, 2, ..., \ell$$
 do
Set $T_j = O(\frac{1}{\alpha^2})$;
 $(\hat{\theta}_j, \hat{\eta}_j) = \operatorname{RegNR}(T_j; \hat{\eta}_1, ..., \hat{\eta}_{j-1})$ (Calling Algorithm 2)
Output: (ℓ, α) -convex lexifair model $\hat{\theta}_\ell$

• At each level *j*, we employ a subroutine in which the Learner plays Online Projected Gradient Descent and the Auditor best responds **ALGORITHM 2:** RegNR: *j*th round **Input:** Number of rounds T, previous estimates $(\eta_1, \ldots, \eta_{i-1})$ Initialize the Learner: $\theta^1 \in \Theta, \eta^1_i \in [0, j \cdot L_M];$ for t = 1, 2, ..., T do Learner plays (θ^t, η^t_i) ; Auditor best responds: $\lambda^t = \lambda_{\text{best}}(\theta^t, \eta_i^t; (\eta_1, \dots, \eta_{j-1}));$ Learner updates its actions using Projected Gradient Descent: $\theta^{t+1} = \operatorname{Proj}_{\Theta} \left(\theta^{t} - \eta \cdot \nabla_{\theta} \mathcal{L}_{i}(\theta^{t}, \eta^{t}_{i}, \lambda^{t}) \right)$ $\eta_i^{t+1} = \operatorname{Proj}_{[0,i:L_M]} \left(\eta_i^t - \eta' \cdot \nabla_{\eta_i} \mathcal{L}_j(\theta^t, \eta_i^t, \lambda^t) \right)$ **Output:** the average play $\hat{\theta} = \frac{1}{\tau} \sum_{t=1}^{T} \theta^t \in \Theta$. and

$$\hat{\eta}_j = \frac{1}{T} \sum_{t=1}^T \eta_j^t \in [0, j \cdot L_M].$$

Auditor plays maximum weight on most violated constraint:

ALGORITHM 3: The Auditor's Best Response (λ_{best}) : *j*th round **Input:** Learner's play (h, η_j) , previous estimates $(\eta_1, \ldots, \eta_{j-1})$ Compute $L_k(h)$ for all groups $k \in [K]$; Find the top *j* elements of vector $(L_1(h), \ldots, L_K(h))$ and call them: $L_{\bar{h}(1)}(h) \ge \ldots \ge L_{\bar{h}(j)}(h)$; **if** $\forall r \le j : L_{\bar{h}(1)}(h) + \ldots + L_{\bar{h}(r)}(h) \le \eta_r$ **then** $\lambda_{out} = 0$; **else** Let $r^* \in \operatorname{argmax}_{r \le j} (L_{\bar{h}(1)}(h) + \ldots + L_{\bar{h}(r)}(h) - \eta_r)$, $\lambda_{out} = \lambda^*$; **Output:** $\lambda_{out} \in \Lambda_j$

- Our ability to prove out of sample bounds crucially relies on our definitional choices that ensure stability.
- Specifically, we show that if:

Our base class H satisfies a standard uniform convergence bound across every group:

For distribution \mathcal{P} and $\delta > 0$ there exists $\beta(\delta)$ such that

$$\Pr_{S}\left[\max_{h\in\mathcal{H},k\in[K]}\left|L_{k}\left(h,S\right)-L_{k}\left(h,\mathcal{P}\right)\right|>\beta(\delta)\right]<\delta$$

 $\textcircled{\sc l}$ We have a model that is approximately convex lexifair on our dataset $S\sim \mathcal{P}^n$

then our model is also appropriately convex lexifair on the underlying distribution.

For every data set *S* sampled *i.i.d.* from \mathcal{P} , if a model *h* satisfies (ℓ, α) -convex lexicographic fairness with respect to *S*, then with probability at least $1 - \delta$ it also satisfies (ℓ, α') -convex lexicographic fairness with respect to \mathcal{P} for $\alpha' = \alpha + 2\ell\beta(\delta)$.

Note that in the case of classification with 0/1 loss, the sample complexity is *polynomial* in the relevant parameters ℓ , α and VC dim.

Suppose \mathcal{H} is a class of binary classifiers with VC dimension $d_{\mathcal{H}}$ For every \mathcal{P} , every data set $S \equiv \{G_k\}_k$ of size *n* sampled *i.i.d.* from \mathcal{P} , if a randomized model $p \in \Delta \mathcal{H}$ satisfies (ℓ, α) -convex lexicographic fairness with respect to *S*, then with probability at least $1 - \delta$ it also satisfies $(\ell, 2\alpha)$ -convex lexicographic fairness with respect to \mathcal{P} provided that

$$\min_{1 \le k \le K} |G_k| = \Omega\left(\frac{l^2 \left(d_{\mathcal{H}} \log\left(n\right) + \log\left(K/\delta\right)\right)}{\alpha^2}\right)$$